# The Qualification of Software Tools in Safety-Related Development

Roberto Bagnara
Prof., University of Parma / CTO, BUGSENG
Member, MISRA C Working Group
Member, ISO JTC1/SC22/WG14 — C std. committee



IWES 2020, February 9th, 2021

### Outline I

- Prologue
- Tool Qualification in General

- 3 Tool Qualification with ISO 26262
- Qualification Kits
- Conclusion

### Acknowledgments

Several slides are courtesy of Marcel Beemster, Solid Sands

### Do You Want to Reason in Assembly?

```
f:
                                 andl
                                         -20(\% \text{rbp}), %eax
.LFB0:
                                 movl
                                         %eax. %eax
  .cfi_startproc
                                addg %rax, -8(\%rbp)
                                 addl $1, -12(\% \text{rbp})
  pushq %rbp
  .cfi_def_cfa_offset 16 .L2:
  .cfi_offset 6, -16
                                movl -12(\% \text{rbp}), \% \text{eax}
  movq
       %rsp, %rbp
                                cmpl -20(\% \text{rbp}), \% \text{eax}
  .cfi_def_cfa_register 6
                               ib .L3
       %edi, -20(\%rbp)
  movl
                              movq -8(\%rbp), \%rax
  movq \$0, -8(\% \text{rbp})
                              popq %rbp
  movl \$0, -12(\% \text{rbp})
                                .cfi_def_cfa 7, 8
  jmp
           .L2
                                 ret
.13:
                                 .cfi_endproc
  movl -12(\% \text{rbp}), \% \text{eax}
```

### Or in C?

```
#include <stdint.h>

uint64_t f(uint32_t n) {
   uint64_t total = 0;
   for (uint32_t i = 0; i < n; ++i) {
     total += i & n;
   }
   return total;
}</pre>
```

# Programming Critical, Resource-Constrained Embedded Systems

C usage is pushed by very strong economic reasons

Unrestricted C has also very serious problems: non-definite behaviors

Ada, too, has non-definite behaviors

Other more defined high-level languages are not portable, flexible or efficient enough

Only two sensible options remain:

- Stick to MISRA C/C++, compile C/C++ to assembly, and reason about programs at the source code level
- Reason about programs at the assembly code level

### Tools Are Badly Needed

Manually checking for MISRA C/C++ compliance is unpractical

Manually translating C/C++ to assembly is unpractical

Tools are needed, for these and for many other activities related to the development of embedded systems

To what extent can the tools be trusted?

### Tool Qualification

The development of safety-critical software is regulated by standards such as:

CENELEC EN 50128 for railway

RTCA DO-178C for airborne software

ECSS-Q-ST-80C for European space applications

IEC 61508 for industry in general

IEC 62304 for medical devices

ISO 26262 for automotive

Due to the complexity of software, development and verification activity, de facto, have to rely on the use of tools

Malfunction of the tools may compromise the integrity of, or fail to detect defects in, the application software

In order to mitigate this risk, the standards prescribe integrity requirement on tools: this is usually called tool qualification

#### CENELEC EN 50128:2011

When tools are being used as a replacement for manual operations, the evidence of the integrity of tools outputs can be adduced by the same process steps as if the output was done in manual operation. These process steps might be replaced by alternative methods if an argumentation on the integrity of tools output is given and the integrity level of the software is not decreased by the replacement.

#### ISO 26262:2018

[The objective of the] qualification of the software tool [is] to create evidence that the software tool is suitable to be used to support the activities or tasks required by the ISO 26262 series of standards (i.e. the user can rely on the correct functioning of a software tool for those activities or tasks required by the ISO 26262 series of standards).

In the different standards, to a varying degree, tools are categorized depending on:

- potential impact of tool failure
- likeliness such failure is detected

Depending on the tool categorization, standards:

- put requirements on tool development
- put requirements on tool documentation
- put requirements on user skills: all software team members including tool users
- put requirements on tool qualification methods

#### An important distinction is between

- tools that can introduce defects in the application software,
   e.g., a C/C++ compiler
- tools that can fail to detect defects in the application software, e.g., a MISRA C compliance checker

#### Another crucial aspect is the scope of use of the tool!

- if the assembly code generated by the C compiler is manually verified, the qualification requirements on the compiler can be softened or eliminated
- if the MISRA C compliance checker is used to justify the elimination of testing activities its qualification requirements are increased

Finally, it must be taken into account that tools qualification can only be performed in the specific context of their actual use

The tool vendor can (and, in some cases, must) supply material that simplifies/enables the tool user to qualify the tool in the specific use context

However, the final responsibility of the tool choice and qualification lies with the tool user

As a consequence, all bragging about "certified tools" amounts to pure and simple marketing hoax

### Tool Qualification with ISO 26262

ISO 26262:2018, Part 8, Section 11: "Confidence in the use of software tools"

Describes the process of tool qualification for a specific use case The qualification process comprises:

- Planning of usage
- Evaluation
- Qualification methods
- Validation and mitigating actions
- Documentation and review

# 1. Planning of usage

#### Determine:

- a) tool identification
- b) tool configuration
- c) tool use case
- d) tool execution environment
- e) maximum ASIL
- f) qualification methods

### 2. Tool evaluation

The Tool Confidence Level (TCL) classes represent the required degrees of confidence in a software tool so that it can be used in a given tool chain, for a given use case, on a given operational environment

#### The classes are:

TCL1 low confidence

TCL2 medium confidence

TCL3 high confidence

## 2. Tool evaluation (cont'd)

The Tool error Detection (TD) classes are meant to capture the confidence in deployed measures that are able to prevent and/or detect malfunctions of a software tool and the consequent production of erroneous output, for a given use case, on a given operational environment

#### The classes are:

- TD1 there is a high degree of confidence that malfunctions and the consequent erroneous outputs will be prevented or detected
- TD2 there is a medium degree of confidence that malfunctions and the consequent erroneous outputs will be prevented or detected
- TD3 there is a low or unknown level of confidence that malfunctions and the consequent erroneous outputs will be prevented or detected

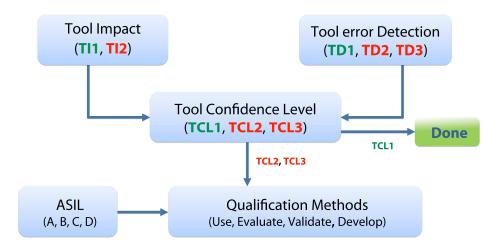
# 2. Tool evaluation (cont'd)

The Tool Impact (TI) classes are meant to capture the possibility that a malfunction of a software tool can introduce or fail to detect an error in a safety-related item or element under development

#### The classes are:

- TI1 the tool can neither introduce nor fail to detect errors
- TI2 the tool can introduce errors and/or fail to detect errors

## 2. Tool evaluation (cont'd)



### 3. Qualification methods

Columns of ISO 26262:2018, Part 8, Section 11, Table 4, TCL3, ASIL C–D:

- (+) Increased confidence from use
- (+) Evaluation of the tool development process
- (++) Validation of the tool
- (++) Development in accordance with a safety standard

### 4. Validation of the tool

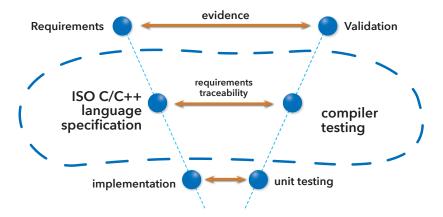
Compliance with the functional specification by testing given the specific use-case

Must take place in the user tool operational environment

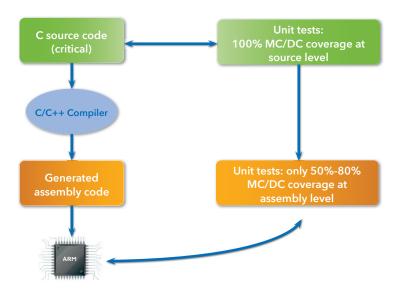
With the precise configuration used in production

Must define mitigations for tool malfunctions

### Compiler Validation



### Compiler Validation: Unit Testing Is Not Enough



# Coverage at Source Level

```
#include < stdint.h>
uint64_t f(uint32_t n)  {
  uint64_t total = 0:
  for (uint32_t i = 0; i < n; ++i) {
    total += i & n;
  return total;
Complete coverage at source level with one test: f(1)
```

### Coverage at Assembly Level: -00

```
-20(\% \text{rbp}), \% \text{eax}
                                   andl
.LFB0:
                                  movl
                                            %eax. %eax
                                  addq
                                            %rax, -8(%rbp)
  .cfi_startproc
  pushq %rbp
                                  addl $1, -12(\% \text{rbp})
  .cfi def cfa offset 16 .L2:
                                  movl -12(\% \text{rbp}), \% \text{eax}
  .cfi_offset 6. -16
  movg %rsp, %rbp
                                  cmpl -20(\% \text{rbp}), \% \text{eax}
  .cfi_def_cfa_register 6
                               ib
                                            .L3
  movl \%edi, -20(\%rbp)
                               movg -8(\% \text{rbp}), \% \text{rax}
  movq \$0, -8(\%rbp)
                              popq %rbp
  movl \$0, -12(\% \text{rbp})
                                  .cfi_def_cfa 7, 8
           .L2
  imp
                                  ret
.L3:
                                   .cfi_endproc
       -12(\% \, \mathsf{rbp}) , \% \, \mathsf{eax}
  movl
```

Without optimization complete coverage also at assembly level with one test: f(1)

### Coverage at Assembly Level: -Ofast

```
f:
                               punpckhdg %mm3, %mm0
                                                            leal
                                                                    2(%rdx), %ecx
IFR0 ·
                               punpckldq %mm3, %mm4
                                                            cmpl
                                                                    %ecx. %edi
                                       %mm4. %mm0
  .cfi_startproc
                               padda
                                                            ibe
                                                                    1.1
  testl
          %edi. %edi
                                       %mm0. %mm1
                                                                    %edi, %ecx
                               paddq
                                                            andl
                                       %edx. %eax
                                                                    $3. %edx
  iе
          17
                               cmpl
                                                            addl
  leal
          -1(%rdi), %eax
                               ib
                                       14
                                                            adda
                                                                    %rcx. %rax
          $3, %eax
                                       %mm1. %mm0
                                                                    %edx, %edi
  cmpl
                               movdga
                                                            cmpl
  ibe
          .L8
                               movl
                                       %edi, %edx
                                                            ibe
                                                                    .L1
          %edi, %edx
                                       $8. %mm0
                                                                    %edx. %edi
  movl
                               psrlda
                                                            andl
          .LC0(%rip), %mm2
  movdaa
                               andl
                                       $-4. \%edx
                                                            adda
                                                                    %rdi, %rax
          %edi, %xmm7
  movd
                               paddq
                                       %mm0. %mm1
                                                            ret
          %eax, %eax
                                       %mm1. %rax
  xorl
                               mova
                                                          17.
  shrl
          $2. %edx
                                       %edx. %edi
                                                            xorl
                                                                    %eax. %eax
                               cmpl
          .LC1(%rip), %mm5
                                       .L11
                                                          .L1:
  movdga
                               iе
  pxor
          %mm1. %mm1
                             13.
                                                            ret
          %mm3, %mm3
  pxor
                               movl
                                       %edi. %ecx
                                                          111 -
          $0, %mm7, %mm6
                                       %edx, %ecx
  pshufd
                               andl
                                                            ret
.L4:
                                       %rcx, %rax
                                                          .L8:
                               addq
          %mm2, %mm0
                                                                    %edx. %edx
  movdga
                               leal
                                       1(%rdx). %ecx
                                                            xorl
                                                                    %eax, %eax
  lb b a
          $1. %eax
                                       %edi. %ecx
                                                            xorl
                               cmpl
          %mm5. %mm2
  bbbsq
                               inb
                                       .L1
                                                           jmp
                                                                    .L3
  pand
          %mm6. %mm0
                               andl
                                       %edi. %ecx
                                                            .cfi_endproc
  movdga
         %mm0. %mm4
                               adda
                                       %rcx. %rax
```

With optimization, f(1) coverage is incomplete at assembly level

### Uses of Static Analyzers in ISO 26262 (Part 6)

# Table 1 — Topics to be covered by modelling and coding guidelines

Topics			ECLAIR			
		Α	В	C	D	ECLAIR
1a	Enforcement of low complexity	++	++	++	++	✓
1b	Use of language subsets	++	++	++	++	✓
1c	Enforcement of strong typing	++	++	++	++	<b>√</b>
1d	Use of defensive implementation techniques	+	+	++	++	✓
1e	Use of well-trusted design principles	+	+	++	++	✓
1f	Use of unambiguous graphical representation	+	++	++	++	-
1g	Use of style guides	+	++	++	++	✓
1h	Use of naming conventions	++	++	++	++	<b>√</b>
1i	Concurrency aspects	+	+	+	+	_

## Uses of Static Analyzers in ISO 26262 (Part 6, cont'd)

Table 3 — Principles for software architectural design

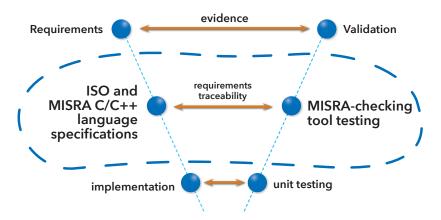
Methods			ECLAIR			
		Α	В	C	D	ECLAIR
1a	Appropriate hierarchical structure of software components	++	++	++	++	<b>✓</b>
1b	Restricted size and complexity of software components	++	++	++	++	<b>✓</b>
1c	Restricted size of interfaces	+	+	+	++	✓
1d	Strong cohesion within each software component	+	++	++	++	✓
1e	Loose coupling between software components	+	++	++	++	<b>√</b>
1f	Appropriate scheduling properties	++	++	++	++	_
1g	Restricted use of interrupts	+	+	+	++	-
1h	Appropriate spatial isolation of the software components	+	+	+	++	_
1i	Appropriate management of shared resources	++	++	++	++	<b>√</b>

## Uses of Static Analyzers in ISO 26262 (Part 6, cont'd)

Table 6 — Design principles for software unit design and implementation

Methods			ECLAIR			
		Α	В	C	D	ECLAIR
1a	One entry and one exit point in subprograms and functions	++	++	++	++	✓
1b	No dynamic objects or variables, or else online test during their	+	++	++	++	✓
	creation					
1c	Initialization of variables	++	++	++	++	✓
1d	No multiple use of variable names	++	++	++	++	✓
1e	Avoid global variables or else justify their usage	+	+	++	++	✓
1f	Limited use of pointers	+	++	++	++	✓
1g	No implicit type conversions	+	++	++	++	✓
1h	No hidden data flow or control flow	+	++	++	++	✓
1i	No unconditional jumps	++	++	++	++	✓
1j	No recursions	+	+	++	++	✓

### MISRA Static Analyzer Validation



# MISRA/HIS Checker Configuration: C is a Large Family of Languages

In C99, there are 112 implementation-defined behaviors

As each i.d.b. can be defined in 2 or more ways, there are more than  $2^{112}\approx 5\times 10^{33}$  possible languages

Actually, choosing integer and floating-types in  $\{8, 16, 32, 64\}$  brings us to more than  $10^{36}$  possible languages (dialects of C)

Alexander's star:  $7.24 \times 10^{34}$  different positions



# C is a Large Family of Languages (cont'd)

Generally speaking, a given compiler can implement, via options, several such dialects of C

For an extreme case, GCC/x86\_64 implements, via options, millions/billions of dialects of C

As a consequences, the tool must adapt to the particular dialect implemented by that compiler with that set of options (possibly for each translation unit)

Further consequence: changing even one compilation option may have important consequences, including analyzing the wrong code!

### 5. Documentation and review

- Software tool criteria evaluation report
- Software tool qualification report, typically resulting in updates to the tool safety manual
- Confirmation review by an independent party

### Qualification Kits

I they are well done, they can decrease the effort of tool qualification by one to two orders of magnitude

#### They must contain:

- Documentation and documentation templates: if a tool safety manual is not there it is a bad sign
- Validated test suites: this requires thousands of tests for a MISRA checker, and tens of thousands of tests for a compiler
- Possibility for users to add their own test cases
- Test automation machinery supplied in source form for inspection
- Possibility of repeating each test completely independently from the qualification kit

### So-called Tool Certificates

A tool cannot be certified, it can be qualified

Marketing people will write just anything to pretend the contrary, e.g.:

- "Usable in [...] ISO 26262 up to ASIL D, TCL1 can be reached"
- "The tool is certified ISO 9001"

In the unfortunate case you end up in court, this kind of things will be demolished in 30 seconds by expert witnesses

Some certification agencies have responsibility for this malpractice

Always do examine the full reports that are an integral part of the so-called certificates

### Conclusion

Tool qualification is an essential requirement for using tools in safety-related developments

We have covered the basic process for ISO 26262, which is not very different from the process described in other functional safety standards

It is a complex process if done in isolation, it is straightforward if done with the help of a good qualification kit

There are advantages besides checking the box:

- Sleeping better (not so with a so-called "tool certificate")
- Decouple application development from tool testing
- Reduced time-to-market

With a good qualification kit, or at least a good validation suite, it is not rocket science

### The End

