

Empirical Evaluation of the Resilience of Novel Non-Algebraic AES S-Boxes to Power Side-Channel Attacks

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Acknowledgments

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Outline

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- (Power) Side-Channel Analysis
 - Definition and goals of the attack
- Countermeasures against SCA
 - Main rationale, issues of existing solutions
 - Why S-Boxes?
- Our contribution
 - Objectives and methodology
 - Results and conclusions

Outline

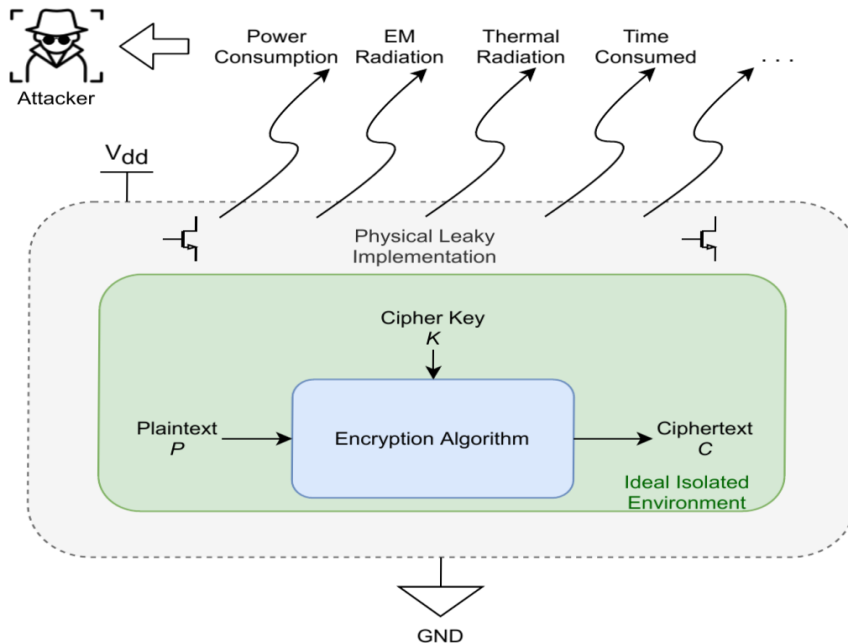
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Side-Channel Analysis

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- **Goal:**
 - Retrieve sensitive data (secret keys)
- **Targets every operating device:**
 - IoT devices, embedded systems
- **Highly effective:**
 - Breaking AES-128:
 - Brute force: 2^{128} → SCA: $2^8 * 2^4 = 2^{12}$
 - Breaking AES-256:
 - Brute force: 2^{256} → SCA: $2^8 * 2^5 = 2^{13}$



Power Analysis

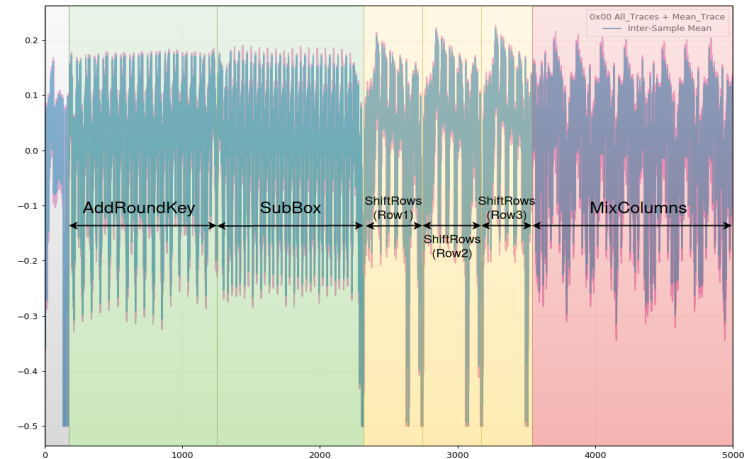
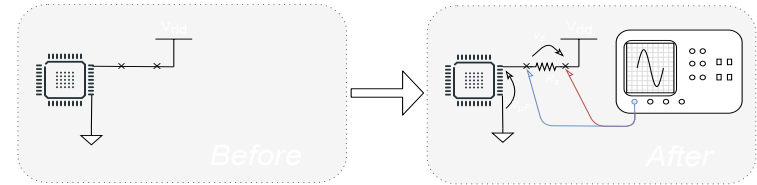
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- **Simple Power Analysis (SPA)**

- A single power trace may suffice
- Human visual analysis
- Reverse Engineering and Timing Attacks

- **Differential/Correlation Power Analysis (DPA/CPA)**

- Multiple power traces are needed
- Based on automated statistical computations
- Correlates the power consumption of the device to the (intermediate) encrypted data (plaintext XOR key)



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SCA Countermeasures

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- **Physical emissions cannot be prevented, but...**
 - 1) Leakage can be reduced
 - 2) What cannot be prevented can be made “unreadable”
- **Countermeasures can be inserted at different levels:**
 - Device-level
 - Balanced transitions, masking, noise circuitry
 - **High impact on area, power consumption and performance**
 - Cryptographic-level
 - Device-independent
 - **New S-Boxes, designed from the ground up to be (hopefully) SCA-resistant**

SCA Countermeasures - Why S-Boxes?

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- **Why target Substitution Boxes?**
 - Crucial component for many block ciphers (AES)
 - **SW implementations:**
 - Look-Up-Tables (LUTs)
 - No memory impact (just replace the original S-Box)
 - Device-independence (one design to rule them all?)
 - Automated heuristic design is possible
 - **Issues:**
 - Difficult to apply on HW implementations
 - Algebraic representation instead of LUTs
 - Adoption
 - Re-standardization, can take years

```
sbox = [  
# 0      1      2      3      4      5      6      7      8      9      a      b      c      d      e      f  
0x63, 0x7c, 0x77, 0x7b, 0xf2, 0x6b, 0x6f, 0xc5, 0x30, 0x01, 0x67, 0x2b, 0xfe, 0xd7, 0xab, 0x76, # 0  
0xca, 0x82, 0xc9, 0x7d, 0xfa, 0x59, 0x47, 0xf0, 0xad, 0xd4, 0xa2, 0xaf, 0x9c, 0xa4, 0x72, 0xc0, # 1  
0xb7, 0xfd, 0x93, 0x26, 0x36, 0x3f, 0xf7, 0xcc, 0x34, 0xa5, 0xe5, 0xf1, 0x71, 0xd8, 0x31, 0x15, # 2  
0x04, 0xc7, 0x23, 0xc3, 0x18, 0x96, 0x05, 0x9a, 0x07, 0x12, 0x80, 0xe2, 0xeb, 0x27, 0xb2, 0x75, # 3  
0x09, 0x83, 0x2c, 0x1a, 0x1b, 0x6e, 0x5a, 0xa0, 0x52, 0x3b, 0xd6, 0xb3, 0x29, 0xe3, 0x2f, 0x84, # 4  
0x53, 0xd1, 0x00, 0xed, 0x20, 0xfc, 0xb1, 0x5b, 0x6a, 0xcb, 0xbe, 0x39, 0x4a, 0x4c, 0x58, 0xcf, # 5  
0xd0, 0xef, 0xaa, 0xfb, 0x43, 0x4d, 0x33, 0x85, 0x45, 0xf9, 0x02, 0x7f, 0x50, 0x3c, 0x9f, 0xa8, # 6  
0x51, 0xa3, 0x40, 0x8f, 0x92, 0x9d, 0x38, 0xf5, 0xbc, 0xb6, 0xda, 0x21, 0x10, 0xff, 0xf3, 0xd2, # 7  
0xcd, 0x0c, 0x13, 0xec, 0x5f, 0x97, 0x44, 0x17, 0xc4, 0xa7, 0x7e, 0x3d, 0x64, 0x5d, 0x19, 0x73, # 8  
0x60, 0x81, 0x4f, 0xdc, 0x22, 0x2a, 0x90, 0x88, 0x46, 0xee, 0xb8, 0x14, 0xde, 0x5e, 0x0b, 0xdb, # 9  
0xe0, 0x32, 0x3a, 0x0a, 0x49, 0x06, 0x24, 0x5c, 0xc2, 0xd3, 0xac, 0x62, 0x91, 0x95, 0xe4, 0x79, # a  
0xe7, 0xc8, 0x37, 0x6d, 0x8d, 0xd5, 0x4e, 0xa9, 0x6c, 0x56, 0xf4, 0xea, 0x65, 0x7a, 0xae, 0x08, # b  
0xba, 0x78, 0x25, 0x2e, 0x1c, 0xa6, 0xb4, 0xc6, 0xe8, 0xdd, 0x74, 0x1f, 0x4b, 0xbd, 0x8b, 0x8a, # c  
0x70, 0x3e, 0xb5, 0x66, 0x48, 0x03, 0xf6, 0x0e, 0x61, 0x35, 0x57, 0xb9, 0x86, 0xc1, 0x1d, 0x9e, # d  
0xe1, 0xf8, 0x98, 0x11, 0x69, 0xd9, 0x8e, 0x94, 0x9b, 0x1e, 0x87, 0xe9, 0xce, 0x55, 0x28, 0xdf, # e  
0x8c, 0xa1, 0x89, 0x0d, 0xbf, 0xe6, 0x42, 0x68, 0x41, 0x99, 0x2d, 0x0f, 0xb0, 0x54, 0xbb, 0x16 # f  
]
```

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- **Our contribution**
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Our contribution - Rationale

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- **What we observed?**
 - **A plethora of new S-Box designs based on:**
 - Chaotic systems and Heuritic methods
 - (Only) theoretical claims of improved SCA resistance
 - **Lack of empirical tests on real world microcontrollers**
 - **Poor results and hard-to-verify benchmarks**
 - Synchronous or asynchronous sampling?
 - Sampling rate?
 - How many traces were collected?
 - Do I have the same microcontroller?

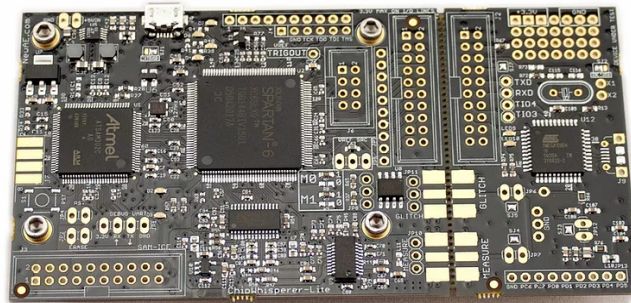
Our contribution – Methodology

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- **What we did?**
 - **Select a subset of the latest S-Box proposals:**
 - #2 LUTs with no resistance to SCA
 - #4 LUTs with (claimed) theoretical resistance to SCA
 - **Empirical analysis leveraging a ChipWhisperer board**
 - 8-bit XMEGA AVR Microcontroller (SW AES-128)
 - Default probe configuration, synchronous sampling
 - Open-source code, completely reproducible
 - **CPA attack, with a Hamming Weight leakage model**

ChipWhisperer-Lite XMEGA

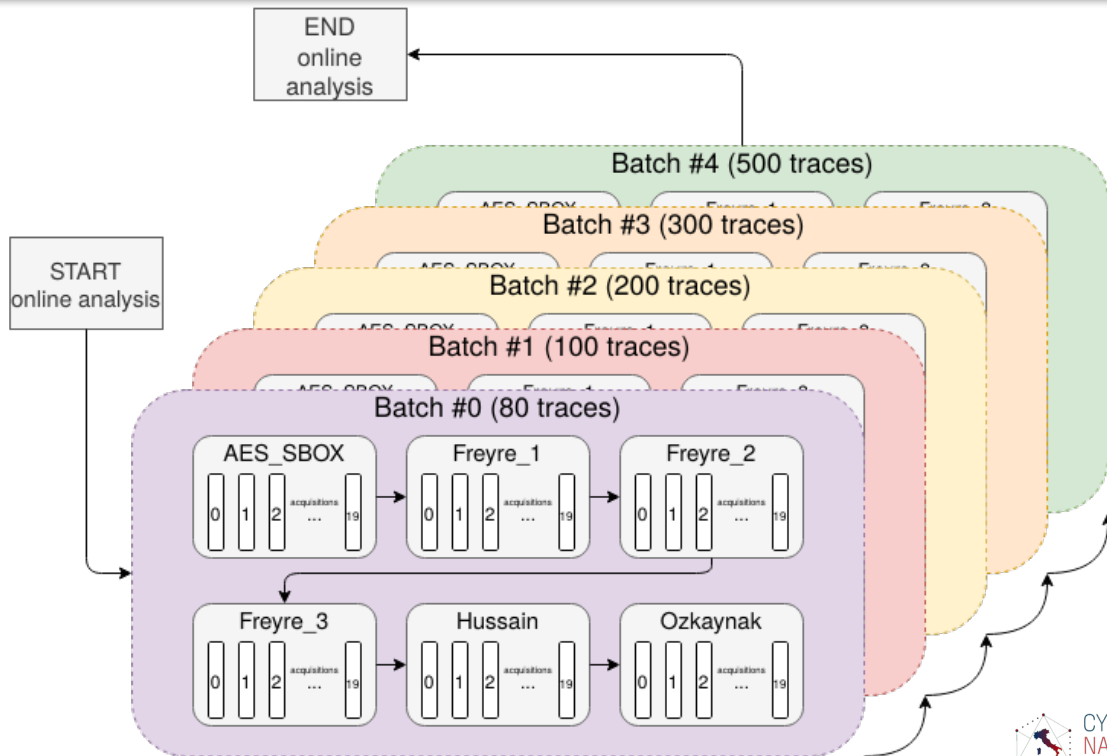
MSRP: \$250 US



Our contribution – Methodology

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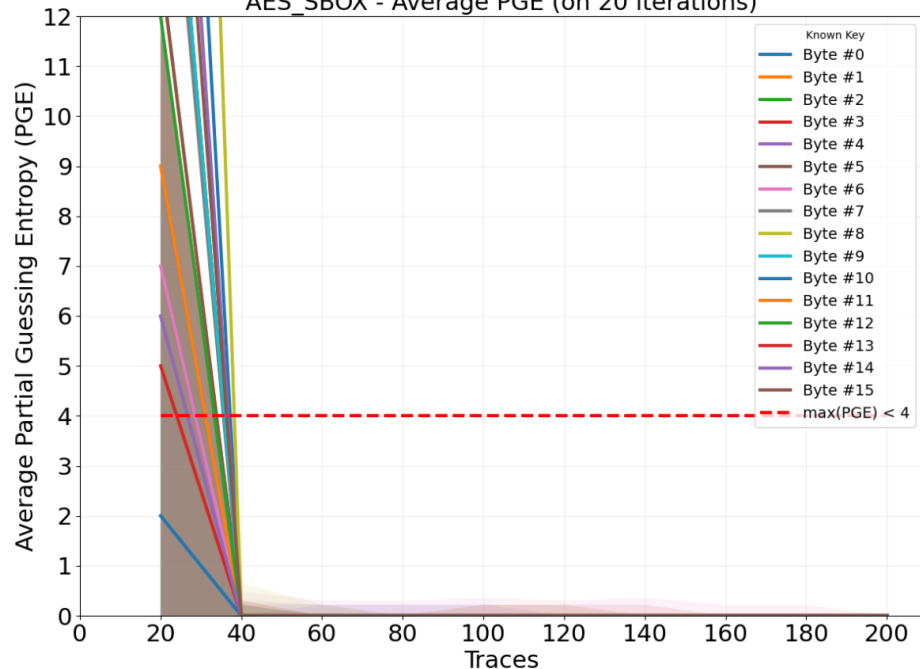
- **Extensive data collection**
 - 20 different datasets are collected for each structure
- **For each S-Box:**
 - 20 CPA attacks are performed
 - Average of the results



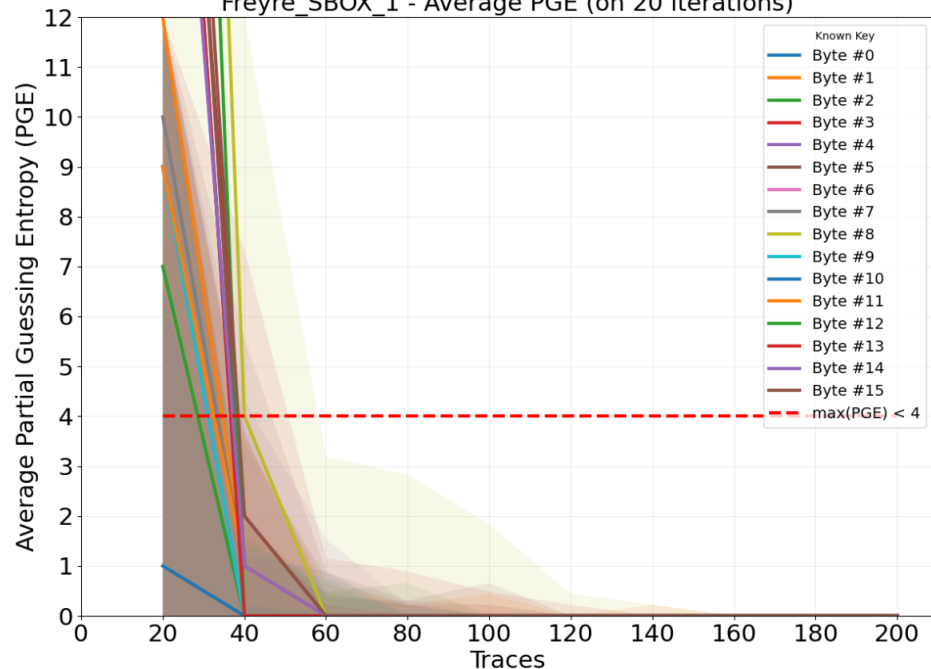
Our contribution – Results

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AES_SBOX - Average PGE (on 20 iterations)



Freyre_SBOX_1 - Average PGE (on 20 iterations)



Our contribution – Observations

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- **What we observed?**
 - **AES, Hussain's and Özkaynak's S-Boxes**
 - Can be broken in less than 50 traces
 - **Freyre's S-Boxes (heuristic methods)**
 - Can be broken, in the best case scenario, within 100 traces
- **Results:**
 - Up to a 2x improvement in the case of Freyre's S-Boxes, seems promising...
 - ... but is it really?

Our contribution – Conclusions

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- **Problem:**
 - **ChipWhisperer-Lite acquires 40 traces per second**
 - Collecting 100 traces only requires 2 seconds!
 - **Sadly, a 2x improvement is not sufficient**
 - Trace capture can be done on-site → time constraints
 - An attack can be carried away later, off-site → no time constraints
- **Conclusions:**
 - Standard benchmarks and reproducible attacks are needed
 - Further improvements on SCA resistance are needed

Thanks for your attention!



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