Failure modes and failures mitigation in GPGPUs: A reference model and its application

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Introduction

General-Purpose Graphics Processing Units (GPGPUs) are now used in fields such as High-Performance Computing and Machine Learning applications

This led researchers and practitioners to focus on the fault-tolerance properties of these devices

We propose a reference model to characterize the failure modes of a GPGPU at different layers







GPGPU Architecture





GPU Architecture – GPU

► GPU – TOP LEVEL:

 global memory (L2 Cache + possible off-chip DRAM), GigaThread Engine (the Global Scheduler), Host Interface, DRAM Channels and the Streaming Multiprocessors (SM).





GPU Architecture Streaming Multiprocessor

- STREAMING MULTIPROCESSOR: Clusters of processors, each processor has its own resources and shares the same L1 Caches
- inside an SM, groups of 32 threads (warps) are scheduled per clock cycle







GPU Architecture Memory Hierarchy

- The memory hierarchy is designed to maximize the bandwidth of data transfers
 - The L2 Cache is designed as a *bandwidth filter* for DRAM













► GPGPUs are made of many similar *replicated* components

- We can group these components in three categories:
 - Computation components (Streaming Multiprocessors, GigaThread Engine and Warp Schedulers, CUDA Cores, Load/Store Units)
 - *Memory components (DRAM, Cache memories, Register Files)*
 - *Communication components (Host Interface, DRAM Channels)*





- Computation Components: Compute service
 - The process of performing arithmetic/logic operations
- Memory Components: *Read, Write* services
- The *read* service is the process of retrieving (read) for a memory location specified by its *address*
- The *write* service is the process of storing (write) data in a memory location specified by its *address*
- Communication Components: Send, Receive services
 - Send is the process of sending data over a channel
- *Receive* is the process of receiving data being transferred over a channel





- ► The GPU can be seen as a *system* providing a *set of services*
- These services are offered by the GPU thanks to the combination of the services offered by its components
- ► The outcome of each service may be *correct* or *incorrect*
- ► To analyze the outcomes of these services, we define two functions:
 - 1. Correct Service: CS(*component*, *service*, *input*, t_i , t_{min} , t_{max}) = (v, t_v)
- 2. Delivered Service: DS(component, service, input, t_i , t_{min} , t_{max}) = (w, t_w)





CORRECT SERVICE

- CS(component, service, input, t_i , t_{min} , t_{max}) = (v, t_v)
- DS(component, service, input, t_i, t_{min}, t_{max}) = (w, t_w)
- a) The service delivered by a component is said to be *correctly valued* if w = v
- b) The service delivered by a component is said to be *correctly timed*, in the presence of a finite timing constraint defined by the lower and upper bounds t_{min} and t_{max} if:

$$t_{\rm i} + t_{\rm min} \le t_{\rm w} \le t_{\rm i} + t_{\rm max}$$





CS(GPU, Compute, x, t_i , t_{min} , t_{max}) = (y, t_y) CORRECT

CONTENT FAILURE

- DS(GPU, Compute, x, t_i , t_{min} , t_{max}) = (z, t_y) WRONG
- DS(GPU, Compute, x, t_i , t_{min} , t_{max}) = (null, t_y) NOT

TIMING FAILURE

- DS(GPU, Compute, x, t_i , t_{min} , t_{max}) = (y, t_z) EARLY $t_z < t_i + t_{min}$ LATE $t_{7} > t_{i} + t_{max}$
- DS(GPU, Compute, x, t_i , t_{min} , t_{max}) = (y, t_z)
- CONTENT AND TIMING FAILURE

- DS(GPU, Compute, x, t_i , t_{min} , t_{max}) = (z, t_z)

ERRATIC $t_{z} < t_{i} + t_{min}$ or $t_{7} > t_{i} + t_{max}$







GPU Failure Detection and Recovery





- The last section of the work consists in a study of the robustness of a GPGPU in three steps:
 - 1. Literature review of the protection techniques developed for GPGPUs (novel techniques or adapted from the CPU domain)
 - 2. What are the *component and/or service* protected by a specific technique
 - 3. What are the *failures* protected by the specific technique





GPU Compute Failure Detection

- Computation components: replication with comparison at different granularities:
- Global replication GPU (adapted by CPU)
- Intra-thread replication Thread level (adapted by CPU)
- Inter-thread replication Warp level (leverages inner redundancy of GPUs)
- Hardware Checkers Comparison is perform using specialized hardware (adapted by CPU, hard for GPU)
- Warp Scheduler protection Software modification technique that allows to detect failures in the warp scheduler units (specific for GPUs)





GPU Memory Failure Detection

Memory components are one of the main failure causes in GPGPUs (high temperatures; frequent accesses to memory to read/write streams of data)

This led NVIDIA to introduce a Single Error Correction Double Error Detection Error Correction Code (SECDED ECC) in all the key memory structures of their GPUs

SECDED ECC can correct 1-bit errors, and detect 2 bits-errors





GPU Communication Failure Detection

GPUs use communication components to communicate with the Host system and the off-chip DRAM, namely the Host Interface and the DRAM Channels

These buffered channels are protected by a Cyclic Redundancy Check (CRC) mechanism with retry

The Host Interface is implemented with the Peripheral Component Interconnect Express (PCIe) standard, which comes with many builtin error-detection capabilities (and CRC with retry)





GPU Computation Failure Recovery

- Transient failures in Computation Components can be recovered with triple replication
 - Perform 3 execution of the same instruction, if there is agreement on at least 2 outputs, the failure can be masked

Permanent failures in SMs require to isolate the failed SM. The proposed approach is to redefine the scheduling strategy to avoid using the corrupted unit, at the cost of reducing throughput by 1/N (where N is the amount of SMs)





GPU Memory Failure Recovery

- The SECDED ECC mechanism always correct 1-bit failures. 2-bits failures are somehow "corrected" with a page (i.e., DRAM row) retirement mechanism, preventing the GPU to use the corrupted memory area in future executions
 - This works both for transient and permanent failures

In the case of a permanent failure in one of the memories, continuously applying this mechanism could reduce the overall throughput of the GPU





GPU Communication Failure Recovery

Transient failures in communication components are corrected by the CRC with retry mechanism

Permanent failures in communication components can not be recovered with the present mechanism. However, it was reported that some permanent failures in the Host Interface could be recovered just by cleaning the GPU connectors





GPU Global Protection

		GPU	Streaming Multiprocessor	Memory Hierarchy	GigaThread Engine	Host Interface	DRAM Channels
	WRONG	Global Replication	Inter/Intra-thread Replication Hardware/Softwa re Checkers	SECDED ECC			
	NOT			Not protected	Not protected	PCI Express Interface CRC with Retry Advanced Error Reporting	CRC With Retry
	EARLY		Inter/Intra-thread Replication	Not protected			
	LATE			Not protected			
RESIL	ERRATIC		Global Replication	SECDED ECC (Value only)			



Graphics Processing Units

Summarizing View





Summarizing View

- The architecture of a GPU is redundant by design, it can be decomposed at different granularities
- GPU and its components provides a set of services which may be correct or incorrect
- Currently available fault-tolerance techniques for GPUs were reviewed with respect to the protected component, and the failure mode

Results showed that some components still need protection

